

Econ 311: Behavioral and Experimental Economics

Prof. Jeffrey Naecker

Wesleyan University

Auctions

What is an Auction?

- ▶ An auction is a way of allocating a single good to one of many interested buyers
- ▶ Buyers typically submit *bids* which can be any positive number
- ▶ Who gets the item, and how much they pay for it, is a function of the all the bids

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- ▶ Note that auctions are a type of game:
 - ▶ Bidders are the players
 - ▶ Bids are the strategies
 - ▶ Payoffs are winnings minus payments

When Are Auctions Used?

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- ▶ Thin markets: unique/rare items such as art, collectibles
- ▶ Thick markets: common/mundane items such as fish, livestock, flowers, online ads
- ▶ Procurement: government seeking private companies to complete large projects
- ▶ Resources: government selling rights oil, land, bandwidth, pollution

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- ▶ Why do economists like auctions?
 - ▶ Efficient allocation of resources

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 - ▶ Winner is person who submits highest bid
 - ▶ Pays either their bid (*first price*) or the next-highest bid (*second price*)
- ▶ Descending bid Dutch auction: start with high price, lower until one person is willing to buy at that price
 - ▶ Called Dutch auction because it's how flowers are sold in Amsterdam

Optimal in Private Values English Ascending-Bid Auction

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- ▶ Thus the only strategy that is not dominated is when you bid your exact valuation

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 - ▶ If you lost, in some cases you could have done better by raising your bid
 - ▶ Thus this bid is dominated by a slightly higher one
- ▶ The only strategy that is not (weakly) dominated is to bid your valuation
- ▶ Note that English auction and second-price sealed bid are *strategic equivalents*: they have same optimal strategy

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- ▶ Optimal strategy is to *shade your bid*

Formal Analysis for First-Price Sealed Bid

- ▶ Suppose you are bidding against a computer which will bid randomly between \$0 and \$10
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 - ▶ Solve to find $b^* = \frac{v}{2}$
- ▶ In general, if N people bidding, best response bid is $b^* = \frac{N-1}{N}v$

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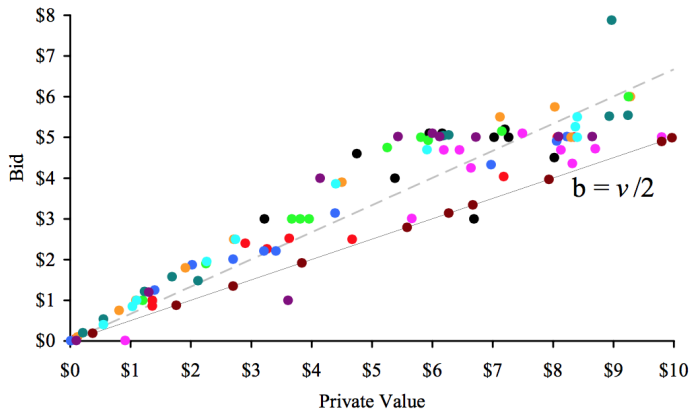
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 - ▶ Can check this payoff is maximized at $b_1 = \frac{v_1}{2}$
 - ▶ Thus player 1 is best-responding
 - ▶ Proof for player 2 is symmetric
 - ▶ Thus this is a NE

Actual Behavior in Private Values First Price Auction



- ▶ Best fit line has slope of about $\frac{2}{3}$ (predicted should be $\frac{1}{2}$)

Optimal Strategy in Private Values Dutch Auction

- ▶ Recall that in Dutch auction, value starts very high and drops until someone stops it
- ▶ Makes sense to pick a point at which you will stop the auction (if someone has not stopped it already)
- ▶ Does it make sense to have stopping point above your value?

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 - ▶ If you win, you will have zero net payoff
 - ▶ If you had waited another second, would have almost certainly won auction at better price
- ▶ Thus again you would want to shade your bid
- ▶ In fact, Dutch Auction and first-price sealed bid are *strategic equivalents*: optimal stopping point in Dutch auction is same as optimal bid in first-price auction

How do People Actually Bid in Auctions?

- ▶ In summary, we found:
 - ▶ In first-price sealed bid and descending clock, bidders should shade their bid (by $\frac{N-1}{N}$)
 - ▶ In second-price sealed bid and English, bidders should bid their valuation (ie no shading of bid)
- ▶ Do people actually play this way?

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- ▶ Do people actually play this way? Let's look at the data from our class exercise

Common Value Auctions

- ▶ Now move to common value auctions
 - ▶ Item is worth the same to everyone
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- ▶ Example: oil lease
 - ▶ Suppose gov't holds land that is worth v to gov't
 - ▶ Gov't decides to auction off land to oil companies
 - ▶ Companies can extract more value from land, say $1.5v$
 - ▶ Companies don't know v exactly; only get signal $v + \epsilon$
 - ▶ Winner is company that submit highest bid
 - ▶ This will be company that gets most optimistic signal
 - ▶ Winner company almost certainly will bid more than land is worth
 - ▶ This is called *winner's curse*